

On Calculations Involving Perturbed Orthogonal Matrices

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Abstract

Many common algorithms employed in Data Science require the calculation of an Orthonormal Basis. Often, this is done using the Gram-Schmidt Algorithm, which has time complexity $O(n^3)$ in the standard case. For low-dimensional spaces, this does not pose a significant computational burden. However, for high-dimensional vector spaces, calculating such a basis can prove unfeasible, so it could be desirable to find faster algorithms that approximate Orthonormal Bases up to some error tolerance ε . In this paper, we examine bounds on error propagation if such an “ ε -almost orthonormal basis” (defined below) is employed.

1. INTRODUCTION

1.1. MOTIVATION

Orthogonal matrices (and their complex counterparts Unitary matrices) are of vital importance in linear algebra, both pure and applied. In a sense, they are the least “disruptive” class of deformations we can apply to a vector space without fixing all of its elements — every linear isometry operator on a finite-dimensional inner product space can be represented by an orthogonal matrix, and conversely, every orthogonal matrix represents an isometry encoded with respect to some particular basis. Thus, every orthogonal matrix Q is necessarily invertible, with $Q^{-1} = Q^T$. This is very nice property, for both theoretical and computational purposes.

From the theoretical perspective, this trivially guarantees the existence of an inverse, and from $1 = \det(I) = \det(QQ^{-1}) = \det(Q)\det(Q^T) = \det(Q)^2$, we get that $\det(Q) = \pm 1$.¹ Additionally, because $Q^{-1}Q = I = QQ^{-1}$, orthogonal matrices encode normal operators, and by the complex spectral theorem they are thus unitarily diagonalizable [Ax15]. In fact, we can guarantee an even stronger property — if Q is orthogonal, then the spectrum $\sigma(Q) \subseteq \{z \in \mathbb{C} \mid |z| = 1\}$. That is, every eigenvalue of Q lies on the unit circle in \mathbb{C} . As such, the group of $n \times n$ orthogonal matrices (denoted $O(n)$) and its subgroups are of great interest from a theoretical perspective.

Orthogonal matrices also have many nice computational properties. Again, let Q be an orthogonal matrix. Then because all eigenvalues have unit modulus, orthogonal matrices cannot magnify errors in computa-

tions [TB97, p. 95]. Furthermore, because $Q^{-1} = Q^T$, inverting Q is essentially trivial, as we can simply swap indexing schemes to yield a semantically faithful representation of Q^{-1} , without having to perform any major calculations. Finally, Orthogonal matrices are of great importance because of their occurrence in important matrix factorizations, such as QR decomposition.

At the time that this paper is being written, there is a large body of ongoing research in numerical linear algebra surrounding “sketching” methods — that is, fast, randomized algorithms for querying large matrices in a way that preserves relevant structure. One classic result in this vein is the *Johnson-Lindenstrauss Lemma* (hence referred to as the JL Lemma), which proves a bounds about the extent to which a family of randomized dimension-reduction algorithms can preserve metric structure. Such tools are necessitated by the ever-growing importance of Big Data and Big Data Analytics techniques in the tech industry.

For our project, we will examine how error resulting from a JL-like dimension reduction scheme can propagate through further computations. To do so, we will re-frame the problem in terms of examining how properties of orthogonal matrices are affected when we apply a small random perturbation matrix. We will examine both additive and multiplicative perturbations, and describe some possible connections to eigenvalue estimation problems.

But first, we will give an overview of any non-standard definitions and notation that we will make use of in our paper. As our randomized algorithms will make some use of concepts from Random Matrix Theory, we will also include an overview of some of the

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¹note that the converse is not true — i.e., for some $n \times n$ matrix M with $\det(M) = \pm 1$, it is not necessarily true that M is orthogonal

relevant concepts in the space below.

1.2. DEFINITIONS AND NOTATIONS

First, some general things: to make things easier for the reader to scan through our paper quickly, we will use \triangle to denote the end of a theorem, definition, remark, or similar, and \blacksquare to denote end of a proof. In general, bases for spaces will be given by \mathcal{B} , while the basis vectors themselves will be denoted \mathbf{e}_i . δ_{ij} will, as usual, refer to the Kronecker delta. In general, we will use the variables i, j, k and ℓ as indices variables, resorting to Greek symbols only if necessary. Unless otherwise stated, \mathcal{U}, \mathcal{V} and \mathcal{W} will refer to finite-dimensional inner product spaces. All other notation (except that described below) should be standard.

DEFINITION 1.1 (Random Matrix Family). We'll use some nonstandard notation for random matrices in our paper. Let \mathcal{A} be a probability distribution with parameters x_1, \dots, x_n , and let $M \in M_{m \times n}(\mathbb{R})$. Then we say $M \sim \mathcal{A}(\mathbf{X}_1, \dots, \mathbf{X}_n)$ iff

$$m_{ij} \sim \mathcal{A}((\mathbf{X}_1)_{ij}, \dots, (\mathbf{X}_n)_{ij}).$$

for all i, j . \triangle

EXAMPLE 1.1.1 (Gaussian random matrix). Let $\boldsymbol{\mu} \in M_{m \times n}(\mathbb{R})$, and let $\boldsymbol{\sigma} \in M_{m \times n}(\mathbb{R}^{\geq 0})$. Then we say $\boldsymbol{\xi} \sim \mathcal{N}(\boldsymbol{\mu}, \boldsymbol{\sigma})$ iff

$$\xi_{ij} \sim \mathcal{N}(\mu_{ij}, \sigma_{ij})$$

for all i, j . We will sometimes denote this by

$$\begin{bmatrix} \xi_{1,1} & \dots & \xi_{1,n} \\ \vdots & \ddots & \vdots \\ \xi_{m,1} & \dots & \xi_{m,n} \end{bmatrix} \sim \begin{bmatrix} \mathcal{N}(\mu_{1,1}, \sigma_{1,1}) & \dots & \mathcal{N}(\mu_{1,n}, \sigma_{1,n}) \\ \vdots & \ddots & \vdots \\ \mathcal{N}(\mu_{m,1}, \sigma_{m,1}) & \dots & \mathcal{N}(\mu_{m,n}, \sigma_{m,n}) \end{bmatrix}$$

if it is more convenient to do so. \triangle

Oftentimes, we'll be interested in matrices where all of the entries are i.i.d. Hence, we'll introduce the following notation:

DEFINITION 1.2. Let $c \in \mathbb{R}$. Then we define

$$c_{m \times n} = c \begin{bmatrix} 1 & \dots & 1 \\ \vdots & \ddots & \vdots \\ 1 & \dots & 1 \end{bmatrix} = \begin{bmatrix} c & \dots & c \\ \vdots & \ddots & \vdots \\ c & \dots & c \end{bmatrix}$$

Thus, we can define a $m \times n$ matrix of i.i.d. Gaussian random variables by $M \sim \mathcal{N}(\mu_{m \times n}, \sigma_{m \times n})$ for some $\mu, \sigma \in \mathbb{R}$ (where $\sigma \geq 0$). \triangle

DEFINITION 1.3. Let $M \in M_{m \times n}(\mathbb{R})$. If we have $M_{ij} < \varepsilon$ for all i, j , then we write

$$M \prec \varepsilon$$

and say M is an ε -bounded matrix. We denote the set of all matrices M such that $-\varepsilon \prec M \prec \varepsilon$ by $B_\varepsilon(0_{m \times n})$,

because under the metric induced by the max norm on $M_{m \times n}(\mathbb{R})$, the set of ε -bounded matrices is simply the open ball about the matrix $0_{m \times n}$.

In the case that we want $M_{ij} \leq \varepsilon$, we of course replace \prec with a \leq symbol, and denote the set of all such matrices by $\overline{B}_\varepsilon(0_{m \times n})$. \triangle

DEFINITION 1.4 (Orthogonal Group). The standard notation for the group of orthogonal matrices on \mathbb{R}^n is $O(n)$. However, this looks confusingly similar to the big-O notation $O(n)$. Hence, we will elect to denote the orthogonal group on \mathbb{R}^n by $\text{Orth}(n)$. \triangle

DEFINITION 1.5. Let $n \in \mathbb{N}$, and let $M \in M_{n \times n}(\mathbb{R})$. Then if M can be expressed as the sum of an orthogonal matrix Q and a matrix ξ with $-\varepsilon \prec \xi \prec \varepsilon$, then we call M a ε -almost orthogonal matrix. \triangle

DEFINITION 1.6. Let \mathcal{V} be a finite-dimensional inner product space. Let $\mathcal{B} = \{\mathbf{e}_1, \dots, \mathbf{e}_n\}$ be a basis for \mathcal{V} , and suppose that $\forall \mathbf{e}_i, \mathbf{e}_j \in \mathcal{B}$ with $i \neq j$, we have

$$|\langle \mathbf{e}_i, \mathbf{e}_j \rangle| < \varepsilon.$$

and $\langle \mathbf{e}_i, \mathbf{e}_i \rangle < 1 + \varepsilon$ for every \mathbf{e}_i . Then we call \mathcal{B} an ε -almost orthonormal basis of V . Note that it will often be useful to restate the conditions above as $\forall \mathbf{e}_i, \mathbf{e}_j \in \mathcal{B}$,

$$|\langle \mathbf{e}_i, \mathbf{e}_j \rangle - \delta_{ij}| < \varepsilon$$

which will be particularly relevant when dealing with double summations. \triangle

To contrast orthonormal bases with ε -almost orthonormal bases, we will often use \mathcal{B}_\perp to denote an orthonormal basis, and \mathcal{B}_ε to denote an ε -almost-orthonormal basis.

This should handle any outstanding notational quirks, so we will continue to the main body.

2. RESULTS

Since we'll be building on results from the midterm project, we'll first reacquaint ourselves with the material therein. To avoid complete redundancy of content, we will take the section of this paper as opportunity to improve the quality and clarity of our exposition. In particular, we will add much greater justification and explanation for our proofs, as many were opaque and generally hard to read. Our analysis will be broken into two main parts: first, dealing with ε -almost orthonormal bases, and then later, dealing with ε -almost orthogonal matrices. These two topics are of course related, as we will detail later. The bulk of our interest will be on the latter, but we will include results about the former as well so as to show two different perspectives on the problem.

2.1. ε -ALMOST ORTHONORMAL BASES

Let \mathcal{V} be an n -dimensional inner product space, and let \mathcal{B}_\perp be an orthonormal basis for \mathcal{V} . One nice property of \mathcal{B}_\perp is that $\forall x \in \mathcal{V}$, we have the representation

$$x = \sum_{i=1}^n \langle x, \mathbf{e}_i \rangle \mathbf{e}_i$$

and hence

$$\|x\|_2^2 = \sum_{i=1}^n |\langle x, \mathbf{e}_i \rangle|^2.$$

To help us build intuition for how it “feels” to work with an ε -almost orthonormal basis instead of an orthonormal basis, we will first prove a result about the extent to which the results above fail.

THEOREM 2.1. *Let \mathcal{V} be finite-dimensional inner product space, and $n = \dim(\mathcal{V})$. Let $\varepsilon > 0$ be given, with the constraint that $\varepsilon < \frac{1}{n}$, and let \mathcal{B}_ε be an ε -almost orthonormal basis of \mathcal{V} . Now, let $x \in \mathcal{V}$, and suppose that x has the representation*

$$x = \sum_{i=1}^n c_i \langle x, \mathbf{e}_i \rangle$$

where $c_i \in \mathbb{R}$, and $\mathbf{e}_i \in \mathcal{B}_\varepsilon$. Then we have

$$\sum_{i=1}^n |c_i| \leq \frac{\|x\|_2}{\sqrt{\frac{1}{n} - \varepsilon}} \quad \triangle$$

Proof. First, we will prove a lower bound for $\|x\|_2$ in terms of the c_i , and use this to obtain the desired result. Observe that by the triangle inequality, we have

$$\left| \sum_{i=1}^n c_i^2 \right| - \left| \langle x, x \rangle - \sum_{i=1}^n c_i^2 \right| \leq |\langle x, x \rangle| = \|x\|_2^2.$$

Removing the superfluous absolute values, this is just

$$\sum_{i=1}^n c_i^2 - \left| \langle x, x \rangle - \sum_{i=1}^n c_i^2 \right| \leq \|x\|_2^2.$$

We seek a bound for the left-hand-side in terms of of the c_i . We'll work with the second term. Note that

$$\langle x, x \rangle = \sum_{i=1}^n \sum_{j=1}^n c_i c_j \langle \mathbf{e}_i, \mathbf{e}_j \rangle.$$

Hence we have

$$\begin{aligned} \left| \langle x, x \rangle - \sum_{i=1}^n c_i^2 \right| &= \left| \sum_{i=1}^n \sum_{j=1}^n c_i c_j (\langle \mathbf{e}_i, \mathbf{e}_j \rangle - \delta_{ij}) \right| \\ &\leq \sum_{i=1}^n \sum_{j=1}^n |c_i c_j| |\langle \mathbf{e}_i, \mathbf{e}_j \rangle - \delta_{ij}| \end{aligned}$$

$$= \sum_{i=1}^n \sum_{j=1}^n |c_i c_j| |(\langle \mathbf{e}_i, \mathbf{e}_j \rangle - \delta_{ij})|$$

Recall that \mathcal{B}_ε is an ε -almost orthonormal basis implies that $\forall \mathbf{e}_i, \mathbf{e}_j \in \mathcal{B}_\varepsilon$, we have $|\langle \mathbf{e}_i, \mathbf{e}_j \rangle - \delta_{ij}| < \varepsilon$, and hence we have

$$\begin{aligned} &\leq \sum_{i=1}^n \sum_{j=1}^n |c_i c_j| \varepsilon \\ &= \varepsilon \sum_{i=1}^n \sum_{j=1}^n |c_i c_j| \\ &= \varepsilon \left(\sum_{i=1}^n |c_i| \right)^2. \end{aligned}$$

This gives us the inequality

$$\sum_{i=1}^n c_i^2 - \varepsilon \left(\sum_{i=1}^n |c_i| \right)^2 \leq \sum_{i=1}^n c_i^2 - \left| \langle x, x \rangle - \sum_{i=1}^n c_i^2 \right|, \quad (1)$$

and so we can chain together with our original bound to yield

$$\left(\sum_{i=1}^n c_i^2 \right) - \varepsilon \left(\sum_{i=1}^n |c_i| \right)^2 \leq \|x\|_2^2$$

By Hölder's inequality for sums, we have

$$\begin{aligned} \left(\sum_{i=1}^n |c_i| \right)^2 &= \left(\sum_{i=1}^n c_i \operatorname{sgn}(c_i) \right)^2 \\ &\leq \left(\sum_{i=1}^n c_i^2 \right) \left(\sum_{i=1}^n \operatorname{sgn}(c_i)^2 \right) \\ &= \left(\sum_{i=1}^n c_i^2 \right) \cdot n \end{aligned}$$

whence

$$\frac{1}{n} \left(\sum_{i=1}^n |c_i| \right)^2 \leq \sum_{i=1}^n c_i^2$$

and so we can chain together with (1) to obtain

$$\frac{1}{n} \left(\sum_{i=1}^n |c_i| \right)^2 - \varepsilon \left(\sum_{i=1}^n |c_i| \right)^2 \leq \|x\|_2^2$$

From which we obtain

$$\begin{aligned} \left(\sum_{i=1}^n |c_i| \right)^2 &\leq \frac{\|x\|_2^2}{\frac{1}{n} - \varepsilon} \\ \sum_{i=1}^n |c_i| &\leq \frac{\|x\|_2}{\sqrt{\frac{1}{n} - \varepsilon}} \end{aligned}$$

as desired. ■

As a corollary, we have the the following result:

COROLLARY 2.1.1. *With all variables quantified as above, we have*

$$|\langle x, \mathbf{e}_i \rangle - c_i| \leq \frac{\varepsilon \|x\|_2^2}{\sqrt{\frac{1}{n} - \varepsilon}}.$$

for every $i = 1, \dots, n$. \triangle

Proof. Observe that

$$\langle x, \mathbf{e}_i \rangle = \left\langle \sum_{j=1}^n c_j \mathbf{e}_j, \mathbf{e}_i \right\rangle$$

Thus we have

$$\begin{aligned} |\langle x, \mathbf{e}_i \rangle - c_i| &= \left| \left\langle \sum_{j=1}^n c_j \mathbf{e}_j, \mathbf{e}_i \right\rangle - c_i \right| \\ &= \left| \left(\sum_{j=1}^n c_j \langle \mathbf{e}_j, \mathbf{e}_i \rangle \right) - c_i \right| \\ &= \left| \sum_{j=1}^n c_j (\langle \mathbf{e}_j, \mathbf{e}_i \rangle - \delta_{ij}) \right| \\ &\leq \varepsilon \left| \sum_{j=1}^n c_j \right| \\ &\leq \varepsilon \sum_{j=1}^n |c_j| \\ &\leq \frac{\varepsilon \|x\|_2^2}{\sqrt{\frac{1}{n} - \varepsilon}} \end{aligned}$$

as desired. \blacksquare

REMARK. Note that if $\varepsilon \ll \frac{1}{n}$, then the error is essentially $O(\varepsilon)$.

2.2. ε -ALMOST ORTHOGONAL MATRICES

We now turn our attention to ε -almost orthogonal matrices. We begin by showing the connection to ε -almost orthonormal bases, and then introducing the problem that will be our main focus for this paper.

2.2.1. INTRODUCTION

THEOREM 2.2. *Let \mathcal{V} be a finite-dimensional inner product space with $\dim(\mathcal{V}) = n$. Let $Q \in \text{Orth}(n)$, and $\xi \in B_\varepsilon(0_{n \times n})$. Then let $M = Q + \xi$, i.e.:*

$$M = \begin{bmatrix} q_{11} + \xi_{11} & \cdots & q_{1n} + \xi_{1n} \\ \vdots & \ddots & \vdots \\ q_{n1} + \xi_{n1} & \cdots & q_{nn} + \xi_{nn} \end{bmatrix}$$

Let $\mathbf{v}_i, \mathbf{v}_j$ be the i^{th} and j^{th} columns of M , respectively. Then

$$\langle \mathbf{v}_i, \mathbf{v}_j \rangle \leq 2\sqrt{n}\varepsilon + n\varepsilon^2.$$

Note that this bound is not tight. In particular, the coefficient of $2\sqrt{n}$ on the ε term is somewhat crude and can be lowered. \triangle

Proof. We have

$$\begin{aligned} \langle \mathbf{v}_i, \mathbf{v}_j \rangle &= \sum_{k=1}^n (q_{ki} + \xi_{ki})(q_{kj} + \xi_{kj}) \langle \mathbf{e}_k, \mathbf{e}_k \rangle \\ &= \sum_{k=1}^n q_{ki}q_{kj} + \xi_{ki}q_{kj} + q_{kj}\xi_{ki} + \xi_{ki}\xi_{kj} \\ &= \langle \mathbf{q}_i, \mathbf{q}_j \rangle + \sum_{k=1}^n \xi_{ki}q_{kj} + q_{kj}\xi_{ki} + \xi_{ki}\xi_{kj} \\ &\leq \sum_{k=1}^n \varepsilon(q_{kj} + q_{ki}) + \varepsilon^2 \\ &= \varepsilon \left(\sum_{k=1}^n (q_{kj} + q_{ki}) \right) + n\varepsilon^2 \\ &\leq \varepsilon \left(2 \sum_{k=1}^n q_{kj} \right) + n\varepsilon^2 \\ &\leq \varepsilon \left(2 \sum_{k=1}^n \frac{1}{\sqrt{n}} \right) + n\varepsilon^2 \quad (*) \\ &= 2\sqrt{n}\varepsilon + n\varepsilon^2 \end{aligned}$$

As desired. Note that to obtain (*), we do not require that $q_{kj} \leq \frac{1}{\sqrt{n}}$ for each k — we are bounding the sum, not the individual terms. \blacksquare

From the theorem, we see that the idea of an ε -almost orthogonal matrix is related to that of an ε -almost orthonormal basis, in that the columns of the ε -almost orthogonal matrix are an $O(\varepsilon)$ almost-orthonormal basis.

THEOREM 2.3. *Let $\varepsilon > 0$ be given, and let $\xi \in B_\varepsilon(0_{n \times n})$. Then $\|\xi\|_F \leq \varepsilon n$.* \triangle

Proof. Observe that

$$\begin{aligned} \|\xi\|_F &= \sqrt{\sum_{i,j \leq n} |\xi_{ij}|^2} \\ &\leq \sqrt{\sum_{i,j \leq n} \varepsilon^2} \\ &= \sqrt{\varepsilon^2 n^2} \\ &= \varepsilon n \end{aligned}$$

as desired. \blacksquare

Now, we'll introduce the problem that will be the main focus of our paper today. An important property of orthogonal matrices is that they can be used to perform fast matrix exponentiation for symmetric matrices. That is, let M be a real symmetric matrix. Then by the spectral theorem, M can be diagonalized by $A\Lambda A^T$, where A is an orthogonal matrix. Using this, we can perform rapid exponentiation of M by utilizing the fact $M^n = A\Lambda^n A^T$, which can be computed quite rapidly. However, we can imagine situations in which we have access to an *almost* orthogonal matrix in some matrix exponentiation problem, but not an *orthogonal* matrix. For instance, imagine we have an IMU measuring the body frame of some machine, say, a UAV, encoded as a matrix B . Say there is measurement error $\xi \sim N(\mathbf{0}, \varepsilon_{m \times n})$ in the hardware, such that B is now an ε -*Almost Orthogonal Matrix*. In this situation, what kinds of errors can we expect in the exponentiation process if we treat B like it's orthogonal? That is, if we assume

$$(BDB^T)^n = BD^n B^T$$

What kinds of errors will this introduce into our calculations? We examine this question below, and prove asymptotic bounds for the errors we might see.

2.2.2. DETERMINISTIC MATRICES

We will begin by examining the problem in the restricted context of non-random matrices, as the proofs are much simpler. Here, we will be primarily interested in two error metrics, which we call δ and Δ , as defined below.

DEFINITION 2.1 (Error Metrics for Diagonalization). Let \mathcal{V} be a finite-dimensional inner product space of dimension $n \in \mathbb{N}$, and let $\varepsilon > 0$ be given. Let $Q \in \text{Orth}(n)$ and $\xi \in B_\varepsilon(0_{n \times n})$, and let $M = A + \xi$. Then we define our Diagonalization error metrics by

$$\delta(k) = \|QD^k Q^T - (MDM^T)^k\|_F^2$$

and

$$\Delta(k) = \|MD^k M^T - (MDM^T)^k\|_F^2.$$

We will refer to the stuff inside the norm in the expression for δ as \mathbf{v}_δ , and the stuff inside the norm in the expression for Δ as \mathbf{v}_Δ . \triangle

Before we begin working with these error metrics, it will be useful to prove some technical lemmas to aid in proving our first theorem.

LEMMA 2.4. *Let Λ, D be diagonal $n \times n$ matrices, and let H be an orthogonal matrix. Then*

$$\text{tr}(HD^k H^T \Lambda^k) = \sum_{i,j} h_{ij}^2 \lambda_i^k(D) \lambda_j^k(\Lambda)$$

\triangle

Proof. We'll start from the left-hand-side, first calculating $HD^k H^T$. Note that indexing variables might get a little wacky here, because we have a lot of matrices in the expression. Observe that, by the standard matrix multiplication formula $c_{ij} = \sum_{k=1}^n a_{ik} b_{kj}$, we have

$$\begin{aligned} (D^k H^T)_{ij} &= \sum_{\alpha=1}^n d_{i\alpha} h_{j\alpha} \\ &= \sum_{\alpha=1}^n \lambda_i^k(D) h_{j\alpha} \delta_{i\alpha} \end{aligned}$$

Where $\delta_{i\alpha}$ is the kronecker delta. Thus, applying this rule a second time, we have

$$\begin{aligned} (H(D^k H^T))_{ij} &= \sum_{\beta=1}^n (H)_{i\beta} (D^k H^T)_{\beta j} \\ &= \sum_{\beta=1}^n h_{i\beta} \sum_{\alpha=1}^n \lambda_\beta^k(D) h_{j\alpha} \delta_{\beta\alpha} \end{aligned}$$

And so finally,

$$\begin{aligned} (HD^k H^T \Lambda^k)_{ij} &= \sum_{\gamma=1}^n (HD^k H^T)_{i\gamma} (\Lambda^k)_{\gamma j} \\ &= \sum_{\gamma=1}^n \left(\sum_{\beta=1}^n h_{i\beta} \sum_{\alpha=1}^n \lambda_\beta^k(D) h_{j\alpha} \delta_{\beta\alpha} \right) \lambda_\gamma^k(\Lambda) \delta_{\gamma j} \\ &= \sum_{\gamma=1}^n \sum_{\beta=1}^n \sum_{\alpha=1}^n h_{i\beta} h_{j\alpha} \lambda_\beta^k(D) \lambda_\gamma^k(\Lambda) \delta_{\beta\alpha} \delta_{\gamma j} \end{aligned}$$

Thus, the trace is given by

$$\begin{aligned} \text{tr}(HD^k H^T \Lambda^k) &= \sum_{\ell=1}^n (HD^k H^T \Lambda^k)_{\ell\ell} \\ &= \sum_{\ell=1}^n \sum_{\gamma=1}^n \sum_{\beta=1}^n \sum_{\alpha=1}^n h_{\ell\beta} h_{\gamma\alpha} \lambda_\beta^k(D) \lambda_\gamma^k(\Lambda) \delta_{\beta\alpha} \delta_{\gamma\ell} \\ &= \sum_{\ell=1}^n \sum_{m=1}^n h_{\ell m} h_{\ell m} \lambda_m^k(D) \lambda_\ell^k(\Lambda) \\ &= \sum_{\ell=1}^n \sum_{m=1}^n h_{\ell m}^2 \lambda_m^k(D) \lambda_\ell^k(\Lambda) \end{aligned}$$

as desired. \blacksquare

COROLLARY 2.4.1. *Let Λ, D, H be as defined above. Then $\text{tr}(\Lambda^k HD^k H^T) = \text{tr}(HD^k H^T \Lambda^k)$. \triangle*

Proof. This follows simply from the fact that trace is invariant under cyclic permutations. \blacksquare

Now, we proceed to the first of our main claims.

THEOREM 2.5. *Let \mathcal{V} be a finite-dimensional inner product space with $\dim \mathcal{V} = n \in \mathbb{N}$, and let $\varepsilon > 0$ be given. Let $Q \in \text{Orth}(n)$ and $\xi \in B_\varepsilon(0_{n \times n})$, and let $M = Q + \xi$. Now, let D be an $n \times n$ diagonal matrix, and let $\sigma_0 = \max\{|\sigma_i(D)|\}$, and $\tau_0 = \max\{|\sigma_j(MDM^T)|\}$ be the largest singular values of D and MDM^T , respectively. Then if $\nu = \max\{\sigma_0, \tau_0\}$, then*

$$\delta(k) \sim \alpha \nu^{2k}$$

for almost all cases, where α is a constant that depends on D and MDM^T . We list edge cases as follows:

- (i) If the eigenvalues of D and MDM^T are equal, then $\delta(k) = 0$.
- (ii) If the singular values of D and MDM^T are equal but some of the eigenvalues are not (i.e., the eigenvalues have the same magnitudes but some have different signs), then $\delta(k) = 0$ for even k and $\delta(k) = \alpha c^{2k}$ for odd k .
- (iii) If the multiplicity of ν in $\sigma(D)$, $\sigma(\Lambda)$ is the same, then whenever k is even, $\delta(k) = o(\nu^{2k})$ (i.e., δ is subexponential in k).

△

Proof. Note that MDM^T is symmetric:

$$\begin{aligned} (MDM^T) &= (M^T)^T D^T M^T \\ &= MDM^T \end{aligned}$$

Thus, by the real spectral theorem, MDM^T is orthogonally diagonalizable, and so there exists an orthogonal matrix P and a diagonal matrix Λ such that $MDM^T = PAP^T$. Hence,

$$\begin{aligned} \delta(k) &= \|QD^k Q^T - (MDM^T)^k\|_F^2 \\ &= \|QD^k Q^T - (PAP^T)^k\|_F^2 \\ &= \|QD^k Q^T - P\Lambda^k P^T\|_F^2 \end{aligned}$$

Since the Frobenius Norm is unitary invariant, it follows that

$$\begin{aligned} &= \|P^T(QD^k Q^T - P\Lambda^k P^T)P\|_F^2 \\ &= \|P^T QD^k Q^T P - \Lambda^k\|_F^2 \\ &= \text{tr}\left((P^T QD^k Q^T P - \Lambda^k)^2\right) \end{aligned}$$

For the sake of concision, let $H = P^T Q$. Then we can write this as

$$\begin{aligned} \delta(k) &= \text{tr}\left((HD^k H^T - \Lambda^k)^2\right) \\ &= \text{tr}(HD^{2k} H^T - HD^k H^T \Lambda^k - \Lambda^k HD^k H^T + \Lambda^{2k}) \\ &\text{by the Lemma and its Corollary, we have} \\ &= \text{tr}(HD^{2k} H^T) + \text{tr}(\Lambda^{2k}) - 2\text{tr}(HD^k H^T \Lambda^k) \\ &= \left(\sum_{i=1}^n \lambda_i^{2k}(D) + \lambda_i^{2k}(\Lambda)\right) - 2 \sum_{i,j \leq n} h_{ij}^2 \lambda_i^k(D) \lambda_i^k(\Lambda). \end{aligned}$$

Let $\nu = \max\{\sigma_0, \tau_0\}$. Suppose that for all $i = 1, \dots, i_0 \leq n$, $j = 1, \dots, j_0 \leq n$, we have

$$\nu = |\lambda_i(D)| = |\lambda_j(\Lambda)|,$$

and without loss of generality, suppose that $i_0 \geq j_0$, and none of the other eigenvalues satisfy this equality. Then we have two cases: $j_0 < n$, and $j_0 = n$ (and hence $i_0 = n$ as well).

- (i) Suppose $j_0 < n$. We evaluate the left-hand term in our expression for $\delta(k)$ first. Because ν was chosen to be the max of the singular values, we have

$$\lim_{k \rightarrow \infty} \frac{\sum_{i=1}^n \lambda_i^{2k}(D) + \lambda_i^{2k}(\Lambda)}{\nu^{2k}} = i_0 + j_0.$$

Now, we examine the right-hand term. By similar reasoning,

$$\begin{aligned} \text{RHS} &= \lim_{k \rightarrow \infty} -2 \frac{\sum_{i,j \leq n} h_{ij}^2 \lambda_i^k(D) \lambda_i^k(\Lambda)}{\nu^{2k}} \\ &= -2 \frac{\sum_{i,j \leq j_0} h_{ij}^2 \lambda_i^k(D) \lambda_i^k(\Lambda)}{\nu^{2k}} \\ &= -2 \sum_{i,j \leq j_0} (-1)^{\beta_{ij}} \end{aligned}$$

where β_{ij} is defined by

$$(-1)^{\beta_{ij}} = \frac{\lambda_i^k(D) \lambda_i^k(\Lambda)}{\nu^{2k}}.$$

(This is because for $i, j \leq j_0$, we have $|\lambda_i^k(D) \lambda_i^k(\Lambda)| = \nu^{2k}$, by definition of j_0). Now, because H is orthogonal, the columns of H are orthonormal, hence we obtain the inequality

$$\begin{aligned} \sum_{i,j \leq j_0} (-1)^{\beta_{ij}} h_{ij}^2 &< \sum_{i=1}^n \sum_{j=1}^{j_0} h_{ij}^2 \\ &= \sum_{j=1}^{j_0} \|\mathbf{h}_j\|_2^2 \\ &= j_0 \end{aligned}$$

with the inequality being strict since $j_0 < n$. This yields the lower bound $i_0 + j_0 - 2j_0 = i_0 - j_0$:

$$0 \leq i_0 - j_0 < \lim_{k \rightarrow \infty} \frac{\delta(k)}{\nu^{2k}}$$

we obtain an upper bound when we always have $\beta_{ij} = 1$, yielding

$$\lim_{k \rightarrow \infty} \frac{\delta(k)}{\nu^{2k}} < i_0 + 3j_0$$

thus exists a constant $\alpha > 0$ such that

$$\lim_{k \rightarrow \infty} \frac{\delta(k)}{\nu^{2k}} = \alpha$$

hence

$$\boxed{\delta(k) \sim \alpha \nu^{2k}}$$

as desired.

- (ii) Now, suppose $i_0 = n = j_0$. Then all of the singular values of D and Λ are the same. Suppose that k of the eigenvalues are of the same sign. If k is even, then $\lambda_i^{2k}(D) = \lambda_i^{2k}(\Lambda)$ for all i , and hence $(-1)^{\beta_{ij}} = 1$. It follows that

$$\begin{aligned} \sum_{i,j \leq n} (-1)^{\beta_{ij}} h_{ij}^2 &= \sum_{i,j \leq n} h_{ij}^2 \\ &= n \end{aligned}$$

hence we have

$$\begin{aligned} \delta(k) &= n + n - 2n \\ &= 0. \end{aligned}$$

In fact, in general, $\delta(k)$ is subexponential in k whenever the multiplicity of ν is the same in both $\sigma(D)$ and $\sigma(\Lambda)$.

If k is odd, then we have

$$\sum_{i,j \leq n} (-1)^{\beta_{ij}} < h_{ij}^2 = n,$$

and so

$$\delta(k) = \left(2n - \sum_{i,j \leq n} (-1)^{\beta_{ij}} h_{ij}^2 \right) \nu^{2k} = \alpha \nu^{2k},$$

where

$$\alpha = \left(2n - \sum_{i,j \leq n} (-1)^{\beta_{ij}} h_{ij}^2 \right).$$

■

COROLLARY 2.5.1. *Quantify all variables as above. Then $\alpha \leq 4n$.* \triangle

Proof. Note that $\forall i, j, |2h_{ij}\lambda_i^k(D)\lambda_i^k(\Lambda)| \leq 2h_{ij}^2\nu^{2k}$. Then because $\sum_{i,j \leq n} h_{ij}^2 = n$, we have

$$\begin{aligned} \delta(k) &= \sum_{i=1}^n \lambda_i^{2k}(D) + \lambda_i^{2k}(\Lambda) - 2 \sum_{i,j \leq n} h_{ij}^2 \lambda_i^k(D) \lambda_i^k(\Lambda) \\ &\leq 2n\nu^{2k} + 2n\nu^{2k} \\ &= 4n\nu^{2k} \end{aligned}$$

as desired. \blacksquare

Before deriving some results for Δ , we should observe that nowhere in the proof above did we utilize the condition that $M = Q + \xi$. Surely, we conjecture, we should be able to use this fact to tighten our bounds. Indeed, this is the case. We outline the proof in the theorem below. But first, we have a small lemma.

LEMMA 2.6. *Let $A, B \in M_{n \times n}(\mathbb{R})$. Then AB and BA share the same eigenvalues.* \triangle

Proof. Let v be an eigenvector of AB , with associated eigenvalue λ . Then we have

$$ABv = \lambda v.$$

Left multiplying each side by B , we have

$$\begin{aligned} B(ABv) &= B\lambda v \\ &= \lambda Bv \\ &= BA(Bv) \end{aligned}$$

hence Bv is an eigenvector of BA with associated eigenvalue λ . Thus, every eigenvalue of AB is an eigenvalue of BA . Conversely, let u be an eigenvector of BA , with associated eigenvalue τ . Then

$$\begin{aligned} A(BAu) &= A\tau u \\ &= \tau Au \\ &= AB(Au) \end{aligned}$$

hence Au is an eigenvector of AB with associated eigenvalue τ . Thus, AB and BA have the same spectrum. \blacksquare

COROLLARY 2.6.1. *Let $Q \in \text{Orth}(n)$, and let $M \in M_{n \times n}(\mathbb{R})$. Then M and QMQ^T have the same spectrum.* \triangle

Proof. Take $A = QM$ and $B = Q^T$, then apply the lemma above. \blacksquare

In the theorem below, we will apply Weyl's inequality for perturbations, hence we give the statement below for completeness:

THEOREM 2.7 (Weyl). *Let $H, P \in M_{n \times n}(\mathbb{R})$, and let $M = H + P$. Suppose M has eigenvalues*

$$\mu_1 \geq \mu_2 \geq \cdots \geq \mu_n,$$

H has eigenvalues

$$\eta_1 \geq \eta_2 \geq \cdots \geq \eta_n,$$

and P has eigenvalues

$$\rho_1 \geq \rho_2 \geq \cdots \geq \rho_n.$$

Then if any two of M, H, P are Hermitian, $\forall i = 1, \dots, n$, we have

$$\eta_i + \rho_n \leq \mu_i \leq \eta_i + \rho_1.$$

\triangle

Now the theorem.

THEOREM 2.8. *Let $Q \in \text{Orth}(n)$ and $D \in M_{n \times n}(\mathbb{R})$ be diagonal, and let $\xi \in B_\varepsilon(0_{n \times n})$. Now, let $M = Q + \xi$, and let $\lambda_1, \lambda_2, \dots, \lambda_n$ be the eigenvalues of QDQ^\top , and let $\tau_1, \tau_2, \dots, \tau_n$ be the eigenvalues of MDM^\top . If $\lambda = \max_i(|\lambda_i|)$, then*

$$|\lambda_i - \tau_i| \leq \lambda(2\varepsilon n + \varepsilon^2 n^2)$$

△

Proof. For this proof, we will make use of the spectral norm, hence we first note some properties of the spectra of QDQ^\top and MDM^\top . By the corollary above, note that the eigenvalues of QDQ^\top are simply the entries of D . Similarly, note that the spectrum of MDM^\top is equivalent to the spectrum of $M^\top MD$. That is,

$$\begin{aligned} \sigma(MDM^\top) &= \sigma(M^\top MD) \\ &= \sigma((Q + \xi)^\top(Q + \xi)D) \\ &= \sigma(D + \xi^\top QD + Q^\top \xi D + \xi^\top \xi D) \end{aligned}$$

and thus

$$\begin{aligned} \|QDQ^\top - MDM^\top\|_2 &= \|D - M^\top MD\|_2 \\ &= \|\xi^\top QD + Q^\top \xi D + \xi^\top \xi D\|_2 \end{aligned}$$

Since QDQ^\top and MDM^\top are symmetric (and thus hermitian), we can apply Weyl's inequality by taking $M = QDQ^\top$, $H = MDM^\top$, and $P = QDQ^\top - MDM^\top$ to obtain the following bound:

$$\begin{aligned} |\lambda_i - \tau_i| &\leq \|D - M^\top MD\|_2 \\ &= \|\xi^\top QD + Q^\top \xi D + \xi^\top \xi D\|_2 \end{aligned}$$

by submultiplicativity of the spectral norm, we have

$$\begin{aligned} &\leq \|D\|_2 \|\xi^\top Q + Q^\top \xi + \xi^\top \xi\|_2 \\ &= \lambda \|\xi^\top Q + Q^\top \xi + \xi^\top \xi\|_2 \end{aligned}$$

and so by the triangle inequality,

$$\begin{aligned} &\leq \lambda (\|\xi^\top Q\|_2 + \|Q^\top \xi\|_2 + \|\xi^\top \xi\|_2) \\ &\leq \lambda (\|\xi^\top\|_2 \|Q\|_2 + \|Q^\top\|_2 \|\xi\|_2 + \|\xi^\top\|_2 \|\xi\|_2) \\ &= \lambda (\|\xi^\top\|_2 + \|\xi\|_2 + \|\xi^\top\|_2 \|\xi\|_2) \\ &= \lambda (2\|\xi\|_2 + \|\xi\|_2^2) \\ &\leq \lambda (2\|\xi\|_F + \|\xi\|_F^2) \\ &\leq \lambda (2\varepsilon n + \varepsilon^2 n^2) \end{aligned}$$

as desired. ■

We have the following corollary:

COROLLARY 2.8.1. *Let $Q \in \text{Orth}(n)$, and let $D \in M_{n \times n}(\mathbb{R})$ be diagonal. Suppose $M = Q + \xi$ where $\xi \in B_\varepsilon(0_{n \times n})$. Let $\lambda_1^k, \lambda_2^k, \dots, \lambda_n^k$ be the eigenvalues of*

D^k , and let $\tau_1(k), \tau_2(k), \dots, \tau_n(k)$ be the eigenvalues of $MD^k M^\top$ for $k \in \mathbb{N}$. If $\lambda = \max_i(|\lambda_i|)$, then

$$|\lambda_i^k - \tau_i(k)| \leq \lambda^k (2\varepsilon n + \varepsilon^2 n^2)$$

△

Proof. This follows immediately by replacing D in the previous theorem by D^k . ■

Using this result, we can tighten some of the bounds we presented earlier.

COROLLARY 2.8.2. *Let $Q \in \text{Orth}(n)$ and let $D \in M_{n \times n}(\mathbb{R})$ be diagonal. Suppose $M = Q + \xi$ where $\xi \in B_\varepsilon(0_{n \times n})$. Then for sufficiently small ε ,*

$$\delta(k) \leq 4n(d(1 + 2\varepsilon n + \varepsilon^2 n^2))^{2k}$$

where $d = \max_i(|d_i|)$

△

Proof. In light of Theorem 2.8 we can assume for small enough ε that the largest singular value λ_i in $\sigma(QDQ^\top)$ corresponds to the largest singular value τ_i in $\sigma(MDM^\top)$. Therefore, we get that ν (the same ν in Theorem 2.5) is bounded by

$$d + d(2\varepsilon n + \varepsilon^2 n^2)$$

Applying Corollary 2.5.1 gives us the desired bound. ■

Note that, for all of the arguments above, by the equivalence of matrix norms, the results also hold (up to a constant multiple) under the Frobenius norm. Thus, we can apply them to our original bounds.

COROLLARY 2.8.3. *Let $Q \in \text{Orth}(n)$, and let $\xi \in B_\varepsilon(0_{n \times n})$. As usual, let $M = Q + \xi$, and let $D \in M_{n \times n}(\mathbb{R})$ be diagonal, and let $\lambda = \max_i(|\lambda_i|)$. Then*

$$\left| \sqrt{\Delta(k)} - \sqrt{\delta(k)} \right| \leq \lambda^k (2 + \varepsilon n + \varepsilon^2 n^2). \quad \triangle$$

Proof. Proving this result is now easy. Let $\lambda = \max_i \lambda_i(D)$. By the reverse triangle inequality, we have

$$\begin{aligned} \left| \sqrt{\Delta(k)} - \sqrt{\delta(k)} \right| &= \|\mathbf{v}_\Delta\|_F - \|\mathbf{v}_\delta\|_F \\ &\leq \|\mathbf{v}_\Delta - \mathbf{v}_\delta\| \\ &= \|MD^k M^\top - QD^k Q^\top\| \\ &\leq \lambda^k (2\varepsilon n + \varepsilon^2 n^2) \end{aligned}$$

as desired. ■

We're interested in the bounds on $\Delta(k)$. We'll use a crude approach here to approximate it.

LEMMA 2.9. *Let $M = Q + \xi$, where $Q \in \text{Orth}(n)$, and $\xi \in B_\varepsilon(0_{n \times n})$. Let $D \in M_{n \times n}(\mathbb{R})$ be diagonal. Then for $\varepsilon < 1$, $\|(MDM^\top)^k\|_2 \leq \|D\|_2^k (1 + O(\varepsilon))$*

△

Proof.

$$\begin{aligned}
\left\| \prod_{i=1}^k MDM^T \right\|_2 &= \left\| \prod_{i=1}^k (Q + \xi)D(Q + \xi)^T \right\|_2 \\
&\leq \prod_{i=1}^k \|(Q + \xi)D(Q + \xi)^T\|_2 \\
&\leq \prod_{i=1}^k \|QDQ^T\|_2 + \|\xi DQ^T\|_2 \\
&\quad + \|QD\xi^T\|_2 + \|\xi D\xi^T\|_2 \\
&= \prod_{i=1}^k \|D\|_2 (1 + 2\varepsilon + \varepsilon^2) \\
&= \|D\|_2^k \sum_{i=1}^{2k} \binom{2k}{i} \varepsilon^i \\
&= \|D\|_2^k (1 + O(\varepsilon)),
\end{aligned}$$

As desired. \blacksquare

REMARK. This bound is exceptionally crude. See computational results section for performance analysis.

3. COMPUTATIONAL RESULTS & RANDOMNESS

We made considerable efforts to obtain theoretical answers to the following questions:

- If we were to draw ξ from $\mathcal{N}(0_{n \times n}, \varepsilon_{n \times n})$, how much tighter would our bounds on $\mathbb{E}[\delta(k)], \mathbb{E}[\Delta(k)]$ be than our current bounds on $\delta(k), \Delta(k)$?
- On a related note, if we sample the entries of our random matrix ξ from another distribution, what
- How can we view the Johnson-Lindenstrauss Lemma from the perspective of almost orthogonal matrices?
- How can $n \times k$ (where $k \ll n$) almost-orthogonal matrices be used in
- Do we always have $\delta(k) > \Delta(k)$?
- Is $\delta(k) \sim \Delta(k)$?

However, after spending a lot of time trying to find literature on these topics and/or trying to prove the bounds ourselves from scratch, we realized we simply didn't have adequate machinery yet to try and approach these problems from a theoretical standpoint. As such, we decided to take a more computational approach, as is detailed in the following section.

Proof. (For result 3.1; plots on next page)

3.1. TESTING δ, Δ BOUNDS

To investigate the first question, we created some python code to test the $\mathbb{E}[\delta(k)]$ and $\mathbb{E}[\Delta(k)]$ when $\xi \sim \mathcal{N}(0_{n \times n}, \varepsilon)$ as a function of n, k , and ε . The ε dependence was not terribly interesting, hence we'll focus mainly on the k, n dependence here. Implementation details can be found in the appendix.

The first big hurdle we have to work past is how we can test the *average*-case performance of the algorithm. Generating acceptably-distributed ξ is trivial, since we're given a particular distribution we want them to follow. However, we also want to make sure that our orthogonal matrices are also generated in an evenly-distributed manner, to ensure that we're not skewing the results by our particular choice of Q . We guarantee this as follows:

Data: Number of dimensions, n

Result: $Q \in \text{Orth}(n)$

$M \leftarrow$ construct a random $n \times n$ matrix with uniformly distributed entries;

$Q, R \leftarrow$ apply QR decomposition to M ;

return Q

Algorithm 1: Algorithm for generating our orthogonal matrix Q

Then, we perform the following computational scheme:

Data: Number of dimensions n, σ, k list, tests per k

Result: Plot of $\log(\delta(k)/\Delta(k))$

initialization of δ list, Δ list;

$Q, R \leftarrow$ apply QR decomposition to M ;

return Q **for** $k \leftarrow \in k$ list **do**

 initialize list of $\delta(k), \Delta(k)$ results;

for $0 \leq i \leq \text{tests per } k$ **do**

$Q \leftarrow$ generate random orthogonal matrix Q ;

$\xi \leftarrow$ sample from $\mathcal{N}(0_{n \times n}, \sigma_{n \times n})$;

$D \leftarrow$ generate a random diagonal matrix D ;

 calculate $\delta(k), \Delta(k)$ in the usual way;

 Push to list of $\delta(k), \Delta(k)$ results;

end

 take the average of the $\delta(k), \Delta(k)$ lists and push to the δ, Δ lists;

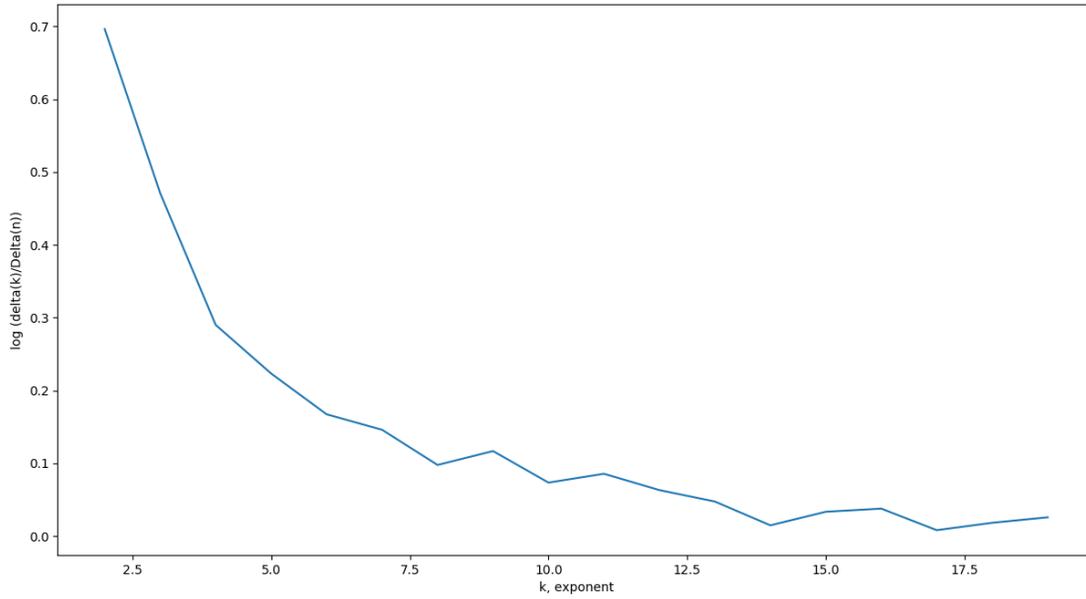
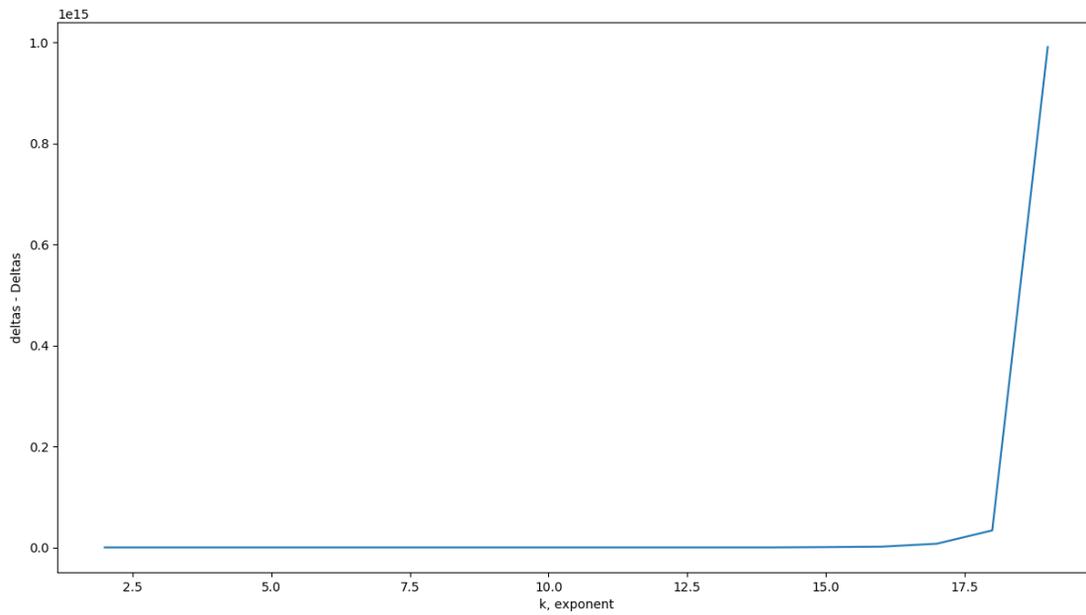
end

perform linear regression on the $\log(\delta), \log(\Delta)$ lists to verify the bound is tight;

Algorithm 2: Algorithm for testing the δ, Δ bounds

Interestingly, upon making plots of $\delta(k) - \Delta(k)$ and $\log\left(\frac{\delta(k)}{\Delta(k)}\right)$, we obtained the following empirical result:

RESULT 3.1. $\delta(k) - \Delta(k) = o(\delta(k)/\Delta(k))$ \triangle

Figure 1: $\log\left(\frac{\delta(k)}{\Delta(k)}\right)$ Figure 2: $\log\left(\frac{\delta(k)}{\Delta(k)}\right)$

hence we see $\log\left(\frac{\delta(k)}{\Delta(k)}\right) \rightarrow 1$, but $\delta(k) - \Delta(k)$ is unbounded, hence we have the desired result. ■

Similarly, observe the following empirical verification of the exponential nature of the bounds:

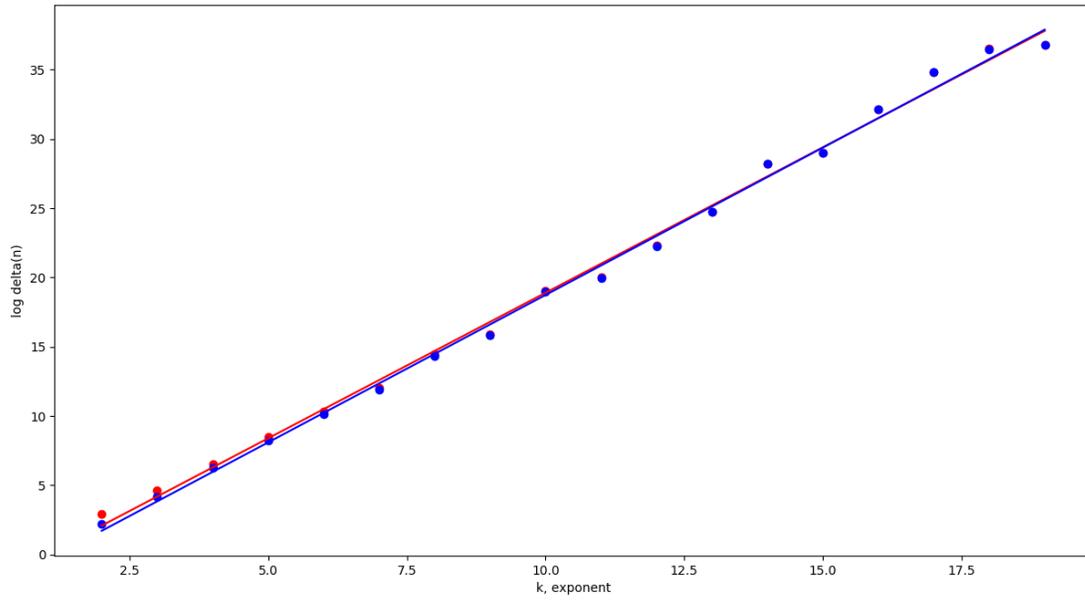


Figure 3: $\log(\delta(k))$

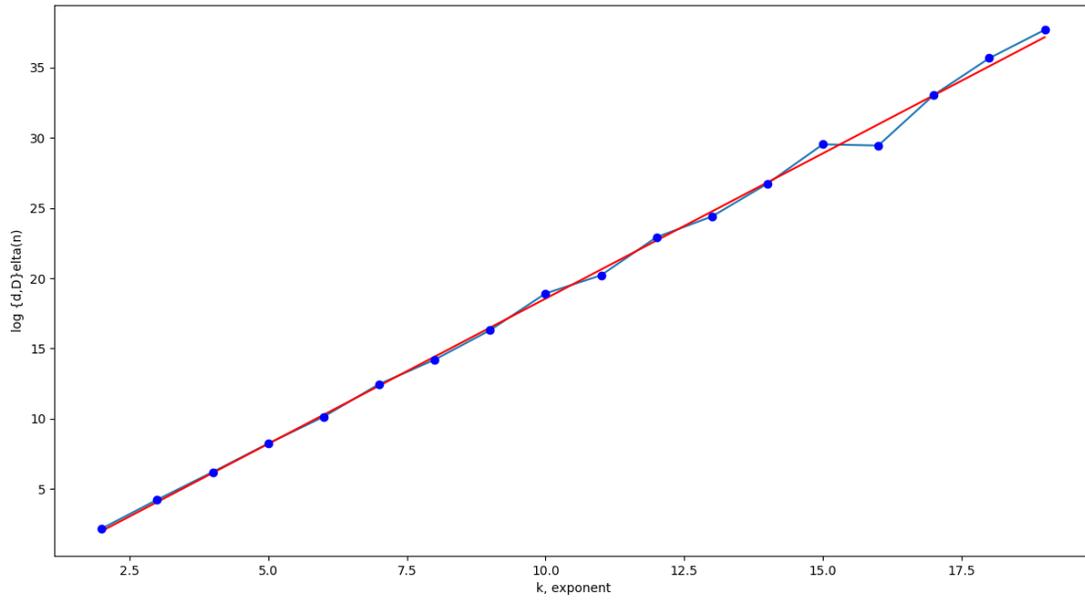


Figure 4: $\log(\Delta(k))$

hence we see $\log\left(\frac{\delta(k)}{\Delta(k)}\right) \rightarrow 1$, but $\delta(k) - \Delta(k)$ is unbounded, hence we have the desired result.

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4. CODE

```

1 import scipy
2 import scipy.linalg
3 # import numpy as np
4 import numpy as np
5 import matplotlib.pyplot as plt
6
7 from tqdm import tqdm
8 from mpl_toolkits import mplot3d
9
10 class Distribution:
11     """
12     Just a simple object wrapper so that we can partially instantiate
13     a distribution for ease of use in defining perturbations below.
14     """
15
16     def __init__(self, distribution, *args, **kwargs):
17         """
18         distribution should be a numpy.random.<distribution_name>
19
20         *args should be whatever parameters the distribution needs to
21         instantiate
22         """
23         self.distribution = distribution
24         self.args = args
25         self.kwargs = kwargs
26         # self.kwargs = kwargs
27
28     def __call__(self, **extra_kwargs):
29         """
30         samples the attached distribution given the args,
31         """
32         return self.distribution(*self.args, **self.kwargs, **extra_kwargs)
33
34
35 # Functions for applying distributions to matrices and stuff
36 def perturb(vec, distribution):
37     """
38     Create a correctly sized perturbation vector and add it to the
39     input vector
40     """
41     return vec + distribution(size=vec.shape)
42
43 def perturb_mat(M, distribution):
44     """
45     Create a correctly sized perturbation matrix and add it to the
46     input matrix
47     """
48     size = M.shape
49     return M + distribution(size = size)
50
51 def construct_almost_basis(n, distribution):
52     """
53     Use a given distribution to perturb basis vectors yielding an
54     (almost) orthonormal basis
55     """
56     I = np.eye(n)
57     for i, row in enumerate(I):
58         I[i] = perturb(row, distribution)
59     return I
60
61 def orthog_error(A):
62     """
63     return error from the square root process as well as the
64     calculation from the "closest" orthogonal matrix
65     """
66     sqrt = scipy.linalg.sqrtm(A.T @ A)
67     invsqrt = np.linalg.inv(sqrt)
68     R = A @ invsqrt
69     return np.linalg.norm(A - R)
70
71 # Helper functions for plot generation
72 def f(x,y):
73     """
74
75

```



```

154         z_label="distance to closest orthogonal matrix")
155
156     plt.show()
157
158     def exp_error(n, A, diag_matrix):
159         """
160         n: exponent to raise to
161
162         Return the error as propogated through when taking the matrix
163         exponential and treating the input
164         """
165         true_matrix = np.linalg.matrix_power(A @ diag_matrix @ A.T, n)
166         est_matrix = A @ (diag_matrix ** n) @ A.T
167         error = np.linalg.norm(true_matrix - est_matrix)
168         return error/(np.linalg.norm(A)**n)
169
170
171     def plot_normal_exp_elbow(min_val=0.0001, max_val=0.5,
172                             num_tests=500, samps_per=50):
173
174         dim = 15
175         sigmas = np.linspace(min_val, max_val, num=num_tests)
176         Z = []
177
178         exps = range(2,70)
179
180         diag_matrix = np.diag(np.random.normal(loc=1.0, scale=.01, size=(dim,)))
181
182         for exp in tqdm(exps):
183             # Reverse the order to start with big errors then go to small
184
185             y_vals = []
186             # y_errs = []
187             for sigma in sigmas:
188                 normal = Distribution(np.random.normal, loc=0.0, scale=sigma)
189                 sub_test = []
190                 for i in range(samps_per):
191                     # print(normal, dims)
192                     almost = construct_almost_basis(dim, normal)
193                     sub_test += [exp_error(exp, almost, diag_matrix)]
194
195                 # Convert to numpy array so we can use nice easy stats without
196                 # writing helper functions
197                 sub_test = np.array(sub_test)
198
199                 # Use the average value
200                 y_vals += [np.mean(sub_test)]
201
202                 # std_dev = np.std(sub_test)
203                 # y_errs += [std_dev/samps_per]
204
205             Z += [y_vals]
206
207     Z = np.array(Z).T
208     X, Y = np.meshgrid(exps, sigmas)
209     # print(X.shape, Y.shape, Z.shape)
210
211     fig = plt.figure()
212     ax = plt.axes(projection='3d')
213     # ax.contour3D(X, Y, Z, 50, cmap='binary')
214     ax.plot_surface(X, Y, np.log(Z), cmap='binary')
215     ax.set_xlabel('exponent')
216     ax.set_ylabel('stard deviation of perturbation')
217     ax.set_zlabel('log normed exponent error')
218     plt.show()
219
220
221
222     def plot_uniform_exp_elbow(min_val=0.0001, max_val=20,
223                               num_tests=100, samps_per=100):
224
225         dim = 15
226         sample_range = np.linspace(min_val, max_val, num=num_tests)
227         Z = []
228
229         exps = range(2,20)
230
231         diag_matrix = np.diag(np.random.normal(loc=1.0, scale=0.5, size=(dim,)))

```

```

232
233 for exp in tqdm(exps):
234     # Reverse the order to start with big errors then go to small
235
236     y_vals = []
237     # y_errs = []
238     for val in sample_range:
239         uniform = Distribution(np.random.uniform, low=-1*val, high=val)
240         sub_test = []
241         for i in range(samps_per):
242             # print(normal, dims)
243             almost = construct_almost_basis(dim, uniform)
244             sub_test += [exp_error(exp, almost, diag_matrix)]
245
246         # Convert to numpy array so we can use nice easy stats without
247         # writing helper functions
248         sub_test = np.array(sub_test)
249
250         # Use the average value
251         y_vals += [np.mean(sub_test)]
252
253         # std_dev = np.std(sub_test)
254         # y_errs += [std_dev/samps_per]
255
256     Z += [y_vals]
257
258 Z = np.array(Z).T
259 X, Y = np.meshgrid(exps, sample_range)
260 # print(X.shape, Y.shape, Z.shape)
261
262 fig = plt.figure()
263 ax = plt.axes(projection='3d')
264 # ax.contour3D(X, Y, Z, 50, cmap='binary')
265 ax.plot_surface(X, Y, np.log(Z+1), cmap='binary')
266 ax.set_xlabel('exponent')
267 ax.set_ylabel('range of uniform perturbation (-y to y)')
268 ax.set_zlabel('log normed (exponent error +1)')
269 plt.show()
270
271
272
273 def plot_uniform_3delbow(min_val=0.001, max_val=10,
274                         num_tests=100, samps_per=5):
275
276     dims = np.array(range(2,30))
277     sample_range = np.linspace(min_val, max_val, num=num_tests)
278     Z = []
279     for dim in tqdm(dims):
280         # Reverse the order to start with big errors then go to small
281
282         y_vals = []
283         # y_errs = []
284         for val in sample_range:
285             uniform = Distribution(np.random.uniform, low=-1*val, high=val)
286             sub_test = []
287             for i in range(samps_per):
288                 # print(normal, dims)
289                 almost = construct_almost_basis(dim, uniform)
290                 sub_test += [orthog_error(almost)]
291
292             # Convert to numpy array so we can use nice easy stats without
293             # writing helper functions
294             sub_test = np.array(sub_test)
295
296             # Use the average value
297             y_vals += [np.mean(sub_test)]
298
299             # std_dev = np.std(sub_test)
300             # y_errs += [std_dev/samps_per]
301
302         Z += [y_vals]
303
304 Z = np.array(Z).T
305 X, Y = np.meshgrid(dims, sample_range)
306 # print(X.shape, Y.shape, Z.shape)
307
308 fig = plt.figure()
309 ax = plt.axes(projection='3d')

```

```

310 ax.contour3D(X, Y, Z, 50, cmap='binary')
311 ax.set_xlabel('dimension of matrix')
312 ax.set_ylabel('range of normal perturbation (-y to y)')
313 ax.set_zlabel('distance to closest orthogonal matrix')
314 plt.show()
315
316 def plot_beta_3delbow(dim=10, min_val=0.001, max_val=10,
317                     num_tests=100, samps_per=10):
318
319     a_range = np.linspace(min_val, max_val, num=num_tests)
320     b_range = a_range.copy()
321
322     Z = []
323     for a in tqdm(a_range):
324         # Reverse the order to start with big errors then go to small
325
326         y_vals = []
327         # y_errs = []
328         for b in b_range:
329             beta = Distribution(np.random.beta, a, b)
330             sub_test = []
331             for i in range(samps_per):
332                 # print(normal, dims)
333                 almost = construct_almost_basis(dim, beta)
334                 sub_test += [orthog_error(almost)]
335
336             # Convert to numpy array so we can use nice easy stats without
337             # writing helper functions
338             sub_test = np.array(sub_test)
339
340             # Use the average value
341             y_vals += [np.mean(sub_test)]
342
343             # std_dev = np.std(sub_test)
344             # y_errs += [std_dev/samps_per]
345
346         Z += [y_vals]
347
348     Z = np.array(Z).T
349     X, Y = np.meshgrid(a_range, b_range)
350     # print(X.shape, Y.shape, Z.shape)
351
352     fig = plt.figure()
353     ax = plt.axes(projection='3d')
354     ax.plot_surface(X, Y, Z, cmap='viridis')
355     ax.set_xlabel('alpha (beta distribution parameter)')
356     ax.set_ylabel('beta (beta distribution parameter)')
357     ax.set_zlabel('distance to closest orthogonal matrix')
358     plt.show()
359
360 def plot_gamma_3delbow(dim=10, min_val=0.001, max_val=10,
361                       num_tests=100, samps_per=10):
362
363     k_range = np.linspace(min_val, max_val, num=num_tests)
364     theta_range = k_range.copy()
365
366     Z = []
367     for k in tqdm(k_range):
368         # Reverse the order to start with big errors then go to small
369
370         y_vals = []
371         # y_errs = []
372         for t in theta_range:
373             gamma = Distribution(np.random.gamma, k, scale=t)
374             sub_test = []
375             for i in range(samps_per):
376                 # print(normal, dims)
377                 almost = construct_almost_basis(dim, gamma)
378                 sub_test += [orthog_error(almost)]
379
380             # Convert to numpy array so we can use nice easy stats without
381             # writing helper functions
382             sub_test = np.array(sub_test)
383
384             # Use the average value
385             y_vals += [np.mean(sub_test)]
386
387             # std_dev = np.std(sub_test)

```

```

388         # y_errs += [std_dev/samps_per]
389
390         Z += [y_vals]
391
392     Z = np.array(Z).T
393     X, Y = np.meshgrid(k_range, theta_range)
394     # print(X.shape, Y.shape, Z.shape)
395
396     fig = plt.figure()
397     ax = plt.axes(projection='3d')
398     ax.plot_surface(X, Y, Z, cmap='viridis')
399     ax.set_xlabel('k (gamma distribution parameter)')
400     ax.set_ylabel('theta (gamma distribution parameter)')
401     ax.set_zlabel('distance to closest orthogonal matrix')
402     plt.show()
403
404
405 def frobenius_tester():
406     dimensions = range(2,100)
407     sigma = .1
408     normal = Distribution(np.random.normal, loc=0.0, scale=sigma)
409
410     tests_per_dim = 10000
411
412     norms = []
413
414     for dim in tqdm(dimensions):
415         this_dim_norms = []
416         for test in range(tests_per_dim):
417             this_dim_norms += [np.linalg.norm(normal(size=(dim,dim)))]
418
419         norms += [np.mean(np.array(this_dim_norms)) - (dim*sigma)]
420
421     print(np.mean(norms))
422     plt.plot(dimensions, norms)
423     plt.xlabel("Dimension of xi matrix")
424     plt.ylabel("Error ||xi||_F relative to (dim * sigma)")
425     plt.title("Error (||xi||_F - dim * sigma); 10000 tests per dim, dim in" +
426             "[2..200]; sigma = 0.1")
427     plt.show()
428
429
430 def frobenius_tester():
431     dimensions = range(2,100)
432     sigma = .1
433     normal = Distribution(np.random.normal, loc=0.0, scale=sigma)
434
435     tests_per_dim = 10000
436
437     norms = []
438
439     for dim in tqdm(dimensions):
440         this_dim_norms = []
441         for test in range(tests_per_dim):
442             this_dim_norms += [np.linalg.norm(normal(size=(dim,dim)))]
443
444         norms += [np.mean(np.array(this_dim_norms)) - (dim*sigma)]
445
446     print(np.mean(norms))
447     plt.plot(dimensions, norms)
448     plt.xlabel("Dimension of xi matrix")
449     plt.ylabel("Error ||xi||_F relative to (dim * sigma)")
450     plt.title("Error (||xi||_F - dim * sigma); 10000 tests per dim, dim in" +
451             "[2..200]; sigma = 0.1")
452     plt.show()
453
454 def delta(n, A, xi, diag_matrix):
455     """
456     n: exponent to raise to
457
458     Return the error as propogated through when taking the matrix
459     exponential and treating the input
460     """
461     M = A + xi
462     true_matrix = np.linalg.matrix_power(M @ diag_matrix @ M.T, n)
463     est_matrix = A @ (diag_matrix ** n) @ A.T
464     error = np.linalg.norm(true_matrix - est_matrix)
465     return error

```

```

466
467 def Delta(n, A, xi, diag_matrix):
468     """
469     n: exponent to raise to
470
471     Return the error as propogated through when taking the matrix
472     exponential and treating the input
473     """
474     M = A + xi
475     true_matrix = np.linalg.matrix_power(M @ diag_matrix @ M.T, n)
476     est_matrix = M @ (diag_matrix ** n) @ M.T
477     error = np.linalg.norm(true_matrix - est_matrix)
478     return error
479
480
481 def get_rms(W, X, y):
482     """ gets the root mean square error """
483     return np.linalg.norm(y - D@W, ord=2)/np.sqrt(D.shape[0])
484
485 def lin_reg(X, y):
486     """ Takes two numpy arrays as inputs, and calculates the coefficients
487     minimizing the least-square error between them """
488     # Solve for optimal weight parameters minimizing regression
489     X = np.column_stack((np.ones_like(X), X))
490     W_opt = np.linalg.solve(X.T @ X, X.T @ y)
491     return W_opt
492
493 def gen_rand_orthog(dim):
494     rand_mat = np.random.rand(dim, dim)
495     q, r = np.linalg.qr(rand_mat)
496     return q
497
498 def test_theorem_4(sigma=.2, plot=True):
499
500     dim = 20
501
502     # dimensions = np.array(range(2,30))
503     exps = range(2,50)
504
505     normal = Distribution(np.random.normal, loc=0.0, scale=sigma)
506
507     tests_per_exp = 100
508
509     deltas = []
510     Deltas = []
511
512     A = gen_rand_orthog(dim)
513     xi = normal(size=(dim, dim))
514
515     for exp in exps:
516         # for exp in tqdm(exps):
517             this_dim_deltas = []
518             this_dim_Deltas = []
519             for test in range(tests_per_exp):
520                 test_mat = np.diag(np.random.normal(loc=1.0, scale=2,
521                                                     size=(dim,)))
522                 this_dim_deltas += [Delta(exp, A, xi, test_mat)]
523                 this_dim_Deltas += [Delta(exp, A, xi, test_mat)]
524
525                 deltas += [np.mean(np.array(this_dim_deltas))]
526                 Deltas += [np.mean(np.array(this_dim_Deltas))]
527
528     deltas = np.array(deltas)
529     Deltas = np.array(Deltas)
530
531     logdeltas = np.log(deltas)
532     logDeltas = np.log(Deltas)
533
534     exps = np.array(exps)
535
536     bd, md = lin_reg(exps, logdeltas)
537     bD, mD = lin_reg(exps, logDeltas)
538
539     if plot:
540         plt.plot(exps, logdeltas, 'ro', exps, logDeltas, 'bo', exps,
541                 (md*exps) + bd, "r", exps, (mD * exps) + bD, "b")
542         plt.xlabel("n, exponent")
543         plt.ylabel("log delta(n)")

```

```

544     plt.show()
545
546     plt.plot(exps, np.log(np.array(Deltas)))
547     plt.plot(exps, logDeltas, "bo")
548     plt.xlabel("n, exponent")
549     plt.ylabel("log {d,D}elta(n)")
550     plt.show()
551
552     plt.plot(exps, np.log(np.log(deltas/Deltas)))
553     plt.xlabel("n, exponent")
554     plt.ylabel("diff of log deltas (n)")
555     plt.show()
556
557     return (bd, md, bD, mD)
558
559 def test_ms():
560     plt.clf()
561     num_tests = 50
562     min_sigma = 0.001
563     max_sigma = 3
564     sigmas = np.linspace(min_sigma, max_sigma, num=num_tests)
565
566     bds = []
567     mds = []
568     bDs = []
569     mDs = []
570
571     for sigma in tqdm(sigmas):
572         bd, md, bD, mD = test_theorem_4(sigma=sigma, plot=False)
573
574         bds += [bd]
575         mds += [md]
576         bDs += [bD]
577         mDs += [mD]
578
579     bds = np.array(bds)
580     mds = np.array(mds)
581     bDs = np.array(bDs)
582     mDs = np.array(mDs)
583
584     print(sigmas.shape, bds.shape, mds.shape, bDs.shape, mDs.shape)
585
586     plt.plot(sigmas, bds, "r", sigmas, mds, "b", sigmas, bDs, "ro",
587             sigmas, mDs, "bo")
588
589     plt.ylabel("slope of best linear fit to log deltas")
590     plt.xlabel("sigma")
591     # plt.savefig("test.png")
592     plt.show()
593
594 def closest_orthog(A):
595     sqrt = scipy.linalg.sqrtm(A.T @ A)
596     invsqrt = np.linalg.inv(sqrt)
597     R = A @ invsqrt
598     return A - R
599
600 def test_orthog_trace():
601
602     numtests = range(100)
603     dims = range(2,50)
604
605     traces = []
606     # trace_errs = []
607
608     traces2 = []
609     # trace_errs2 = []
610
611     for dim in tqdm(dims):
612         this_dim_traces = []
613         this_dim_traces2 = []
614
615         for test in numtests:
616             rand_mat = 2*np.random.randn(dim,dim) + 1
617             this_dim_traces += [np.trace(closest_orthog(rand_mat))]
618
619             Q, _ = np.linalg.qr(rand_mat)
620             this_dim_traces2 += [np.trace(Q)]
621

```

```

622     this_dim_traces = np.array(this_dim_traces)
623     this_dim_traces2 = np.array(this_dim_traces2)
624
625     # sigma = np.std(this_dim_traces)
626     mu = np.mean(this_dim_traces)
627     traces += [mu]
628     # trace_errs += [sigma]
629
630     # sigma2 = np.std(this_dim_traces2)
631     mu2 = np.mean(this_dim_traces2)
632     traces2 += [mu2]
633     # trace_errs2 += [sigma2]
634
635     plt.plot(dims, traces, "b", dims, traces2, "r")
636
637     plt.show()
638
639
640 def plot_beta_3delbow(dim=10, min_val=0.001, max_val=10,
641                      num_tests=100, samps_per=10):
642
643     a_range = np.linspace(min_val, max_val, num=num_tests)
644     b_range = a_range.copy()
645
646     Z = []
647     for a in tqdm(a_range):
648         # Reverse the order to start with big errors then go to small
649
650         y_vals = []
651         # y_errs = []
652         for b in b_range:
653             beta = Distribution(np.random.beta, a, b)
654             sub_test = []
655             for i in range(samps_per):
656                 # print(normal, dims)
657                 almost = construct_almost_basis(dim, beta)
658                 sub_test += [orthog_error(almost)]
659
660             # Convert to numpy array so we can use nice easy stats without
661             # writing helper functions
662             sub_test = np.array(sub_test)
663
664             # Use the average value
665             y_vals += [np.mean(sub_test)]
666
667             # std_dev = np.std(sub_test)
668             # y_errs += [std_dev/samps_per]
669
670         Z += [y_vals]
671
672     Z = np.array(Z).T
673     X, Y = np.meshgrid(a_range, b_range)
674     # print(X.shape, Y.shape, Z.shape)
675
676     fig = plt.figure()
677     ax = plt.axes(projection='3d')
678     ax.plot_surface(X, Y, Z, cmap='viridis')
679     ax.set_xlabel('alpha (beta distribution parameter)')
680     ax.set_ylabel('beta (beta distribution parameter)')
681     ax.set_zlabel('distance to closest orthogonal matrix')
682     plt.show()

```

almost.py